Learning Goals

- Define variance and standard deviation
- State Chebyshev inequality and Chernoff inequality
- Compare the conditions and strengths of Markov, Chebyshev and Chernoff inequalities
- Understand the main idea and steps in the proofs of these bounds
- Intuition for the bounds given by the simplified forms of the Chernoff bound

Chebyshev Inequality

Definition

The *variance* of a random variable *X* is

 $Var[X] := E[(X - E[X])^2] = E[X^2] - (E[X])^2$. Its square root, $\sqrt{Var[X]}$, is the *standard deviation* of X, and is often denoted as σ .

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Theorem (Chebyshev Inequality)

For any $\alpha > 0$, $\Pr[|X - \mathbf{E}[X]| > \alpha \sigma] \leq \frac{1}{\alpha^2}$.

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For any
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, $\Pr[|X - \mathbf{E}[X]| > \alpha \sigma] \le \frac{1}{\alpha^2}$.

Proof.

Apply Markov inequality to the random variable $(X - \mathbf{E}[X])^2$:

$$\Pr[|X - \mathbf{E}[X]| \ge \alpha \sigma] = \Pr[(X - \mathbf{E}[X])^2 \ge \alpha^2 \operatorname{Var}[X]] \le \frac{1}{\alpha^2}.$$



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- \Rightarrow A distribution where $|X \mathbf{E}[X]|$ takes two values: 0 and $\alpha \sigma$
- \Rightarrow *X* takes three values: $\mathbf{E}[X]$, $\mathbf{E}[X] + \alpha \sigma$ and $\mathbf{E}[X] \alpha \sigma$, with equal probability on the latter two values.

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If X and Y are independent random variables, then $E[XY] = E[X] \cdot E[Y]$, and Var[X + Y] = Var[X] + Var[Y].

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Without independence, Var[X + Y] in general is not equal to Var[X] + Var[Y].

Theorem

Let X_1, X_2, \cdots be independently, identically distributed (i.i.d.) random variables, and each has finite variance. For each $n \ge 1$, let \overline{X}_n be $\frac{1}{n} \sum_{i=1}^n X_i$. Then for any $\delta > 0$, $\lim_{n \to \infty} \Pr[|\overline{X}_n - \mathbf{E}[\overline{X}_n]| > \delta] = 0$.

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The right hand side goes to 0 as *n* goes to infinity.



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 - Bernstein and Chernoff exploited the idea by looking at $f(x) = e^{\lambda x}$.

Chernoff Bound: I.I.D. Case

Let X_1, \dots, X_n be i.i.d. Bernoulli variables, such that $\Pr[X_i = 1] = p$ and $\Pr[X_i = 0] = q := 1 - p$ for each i. Define $X = \sum_{i=1}^n X_i$.

Theorem (Chernoff Bound)

For any t > 0,

$$\Pr\left[X > (p+t)n\right] \le \exp\left\{\left(-(p+t)\ln\frac{p+t}{p} - (q-t)\ln\frac{q-t}{q}\right)n\right\}.$$

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Proof of Chernoff Bound (Cont.)

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The same proof yields the same bound for $\Pr[X \leq (p-t)n]$.



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What if X_1, \dots, X_n always take values from [0, 1] but not necessarily $\{0, 1\}$? Suppose $\mathbf{E}[X_i] = p_i$, $q_i = 1 - p_i$ for each i, and let $p = \frac{1}{n} \sum_i p_i$ and q = 1 - p.

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Useful Forms of Chernoff Bound

Corollary

Let X_1, \dots, X_n be independently distributed on [0, 1] and $X = \sum_i X_i$.

• For all t > 0,

$$\Pr[X > E[X] + t], \Pr[X < E[X] - t] \le e^{-2t^2/n};$$

• For any $\epsilon < 1$,

$$\Pr\left[X > (1+\epsilon) \operatorname{\mathbf{E}}\left[X\right]\right] \leq \left(\frac{e^{\epsilon}}{(1+\epsilon)^{1+\epsilon}}\right)^{\operatorname{\mathbf{E}}\left[X\right]} \leq \exp\left(-\frac{\epsilon^{2}}{3} \operatorname{\mathbf{E}}\left[X\right]\right);$$

$$\Pr\left[X < (1 - \epsilon) \operatorname{E}\left[X\right]\right] \le \left(\frac{e^{-\epsilon}}{(1 - \epsilon)^{1 - \epsilon}}\right)^{\operatorname{E}\left[X\right]} \le \exp\left(-\frac{\epsilon^2}{2} \operatorname{E}\left[X\right]\right).$$

Useful Forms of Chernoff Bound (Cont.)

Corollary ((Cont.))

• For any $\epsilon > 1$,

$$\Pr\left[X > (1+\epsilon) \operatorname{E}\left[X\right]\right] \leq \left(\frac{e^{\epsilon}}{(1+\epsilon)^{1+\epsilon}}\right)^{\operatorname{E}\left[X\right]} \leq \exp\left(-\frac{\epsilon}{3} \operatorname{E}\left[X\right]\right);$$

• If $t > 2e \mathbf{E}[X]$, then

$$\Pr[X > t] \le 2^{-t}$$
.



Proof Sketch.

• Let f(t) be $(p+t) \ln \frac{p+t}{p} + (q-t) \ln \frac{q-t}{q}$. Show $f(t) \ge 2t^2$ by showing f(0) = f'(0) = 0 and $f''(t) \ge 4$ for all $0 \le t \le q$ followed by Taylor's theorem with remainder.

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- Let g(x) be f(px), then g'(0) = pf'(px), and so g(0) = g'(0) = 0. Show $g'(1) > p \ln 2 > \frac{2}{3}p$. Deduce that for $x \in (0, 1), g(x) \ge px^2/3$.



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- Set h(x) := g(-x). Then h'(x) = -g'(-x), and h(0) = h'(0) = 0. Show then $h''(x) \le p$ for $x \in (0, 1)$. Deduce that $h(x) \ge px^2/2$.



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See assigned reading for more details. Or take them as exercises.



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