

Dynamic programming; running time

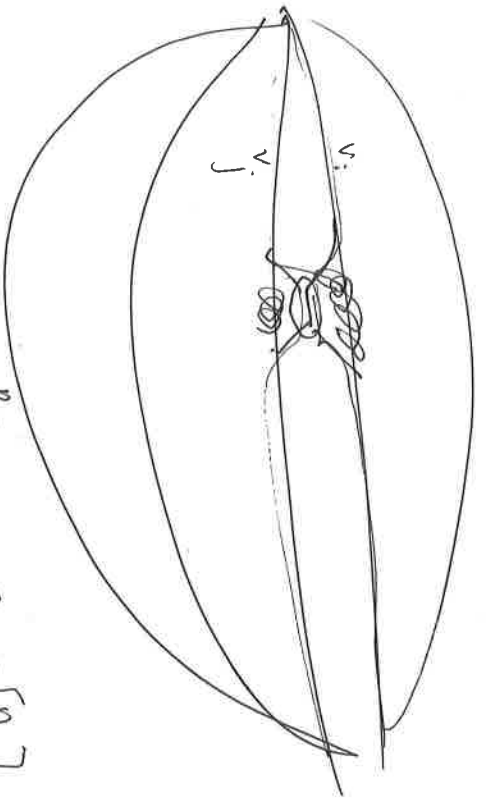
poly  $(B, K, n)$ .

range.  $B_n$ . Input length:  $O(n + \log B + \log K)$

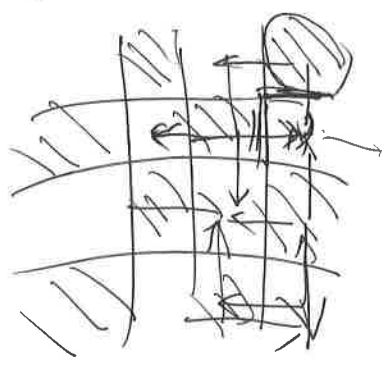
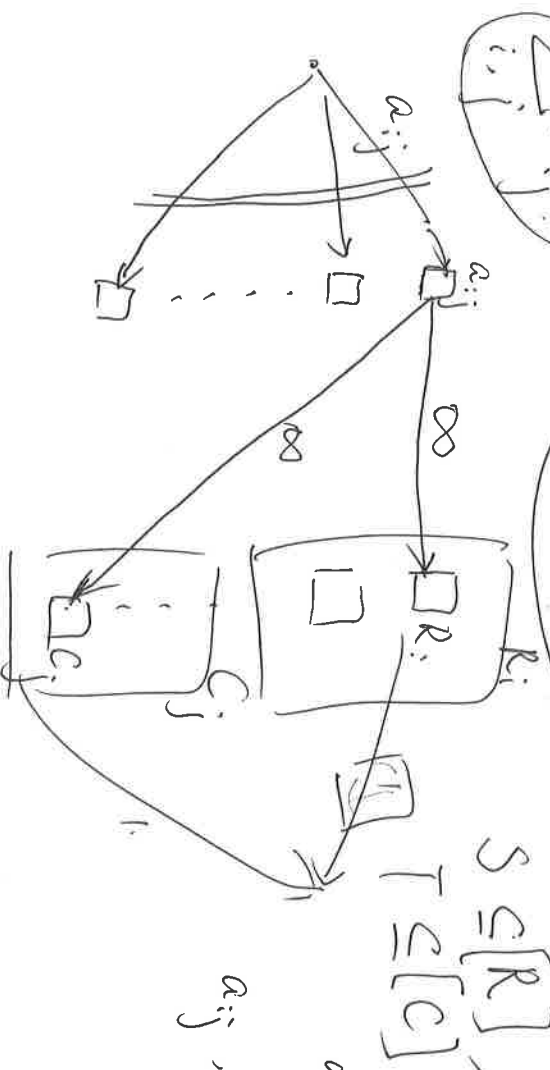
$$B = 2^{\log B}$$

Recall in Subset Sum, the reduction produces a  $K \approx (1+m)^{m+1} \approx 2^{m \log m}$

$$e = (v_i, v_j)$$



$$\sum_{i,j} a_{i,j}$$



$$S \subseteq [R], T \subseteq [C]$$

are on the side of  $t$ ,  $a_{i,j}$ , either  $i \in S$  or  $j \in T$ ,