

Max Flow and Min Cuts: Basics

- Definitions: flow network, flow and its value, s - t cut and its capacity
- Max flow min cut theorem
- Ford-Fulkerson algorithm and its consequences
 - Residual graph, augmenting path; steps of the algorithm
 - Running time for integral capacity networks: $O(mC)$, $C :=$ sum of capacities on edges leaving s .
 - Existence of integral maximum flow on integral capacity networks.
 - Proof of max-flow-min-cut theorem. Computing min cut.
- Existence of polynomial time algorithm (e.g. Edmonds-Karp algorithm)

Max Flow and Min Cuts: Applications

- Bipartite matchings.
 - Maximum bipartite matchings.
 - Hall's theorem and its proof.
- Edge-disjoint paths. (Menger's theorem)
- Dilworth theorem. Partially ordered sets, chains, antichains.
- Vertex cover in bipartite graphs. (König-Egerváry theorem)
- Circulations (possibly with edge lower bounds).
- Image segmentation
- Project selection

NP and computational intractability

- Definitions: Polynomial time reductions, P, NP, NP-complete problems.
- Cook-Levin Theorem: SAT is NP-complete.
- Proving a problem to be NP-complete
 - Prove membership in NP (give a polynomial length certificate and poly-time verifier)
 - Reduce a known NP-complete problem A to the problem B at hand, showing that an instance of A has a solution if and only if the image instance of B does.

Classical NP-complete problems

- 3-SAT
- Packing: Independent Set
- Covering: Vertex Cover, Set Cover
- Sequencing: Hamiltonian Cycle/Path, Traveling Salesman Problem
- Partitioning: 3-Dimensional Matching, Graph Coloring
- Numerical: Subset Sum, Number Partitioning, Knapsack

Exact solutions to NP-hard problems

- Fixed parameter tractable problems: clever exhaustive search.
 - Bounding the running time for recursive algorithms.
 - Common technique: find a way to minimize the “branching” factor.
 - Example: Vertex cover
- Greedy and dynamic programming
 - Independent set on the trees
 - Coloring circular arcs

Approximation algorithms

- Definition: α -approximation algorithms, approximation ratio
- Example greedy approximation algorithms.
 - Load balancing.
 - Center selection (k -center)
 - Set cover
 - Harmonic function $H(n) = \sum_{i=1}^n \frac{1}{i} \approx \log n$.

- Common methods
 - Bounding the optimal
 - Express greedy or termination condition
 - Minimum \leq Average \leq Maximum
 - Summing over conditions and noticing double (or multiple) counting
- (Fully) Polynomial Time Approximation Schemes.
 - Example: Knapsack.
 - Dynamic programming for small instances.
 - Round numbers to get a small instance.

Randomized Algorithms

- Probability events, independence.
- Union bound.
- Random variables, expectations.
 - Example: Geometric distribution: toss a (biased) random coin, expected number of tosses till a Heads vs. number of tosses so that a Heads shows up with high probability.
 - Example: Contention resolution: waiting time for one task and all tasks
- Calculating probability of the intersection of events using conditional probabilities
 - Example: Karger's min cut algorithm

Linearity of Expectations

- Linearity of expectations: independence not required!
 - Method: Decompose the count of a random object and use indicator variables.
- Examples:
 - Guessing cards
 - Coupon collection
 - Quicksort
 - Max SAT: $\frac{8}{7}$ -approximation.